CSCI 2021: Binary, Integers, Arithmetic

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Logisitcs

Reading

Bryant/O'Hallaron Ch 2.1-2.3

Goals

- Finish C Overview
- Binary Notation and Relations
- Integers in binary
- Arithmetic operations

Assignments

- P1 Ongoing
- Lab03 on File Input
- HW03 on Binary Ints

Date		Event
Wed	01-Feb	Structs / File IO
		Lab03 / HW03
Fri	03-Feb	Binary Ints/Chars
Mon	06-Feb	Binary Ints/Chars
Tue	07-Feb	Unified Office Hours
		~ in Lind 325
		Lab03/HW03 Due
Wed	08-Feb	Lec: Practice Exam
		Lab04/HW04: Review
		P1 Due
Fri	10-Feb	**Exam 1**

Prev Lab/HW Solutions Posted

Questions??

Announcements: Unified Office Hours

- To help students finish up Project 1, all staff office hours will be held in Lind 325 on Tue 07-Feb. Staff will likely be present all day in that room with a shared Queue
- Expect these office hours session to be popular (e.g. crowded) and plan on there being some wait time before staff get to you
- Have your Office Hours Help Request Form ready for staff when your turn comes up: `https://z.umn.edu/2021help`

Exam 1 Logistics

- In-person in class on Fri 10-Feb
- Exam runs lecture period: 50min
- Expect 3 sides of paper (front, back, front)

Open Resource Exam

Open Resource Exam Rules

- Sign the log on turning in your exam to show attendence
- Silence your devices and keep screens visible to instructor
- Protect your work from theft
- You may be asked to show your Student ID

Can Use Your Own

Physical, Digital, or Online⋯

- Notes, Slides, Dictionary
- Your own previous Exams
- Textbook(s) (online ok)
- Editor, Compiler, Vole, SSH
- Your code / Instructor code
- Locally stored webpages
- Online Manual Pages
  `http://man.he.net/`
  ex: search for `ascii`

Cannot Use

- General Internet
- Piazza Discussion
- Online calculators, converters, tables
- Chat, Texting, IM, etc.
- Communication with anyone but Instructor/Proctor

If you aren't sure of something, ask

Unsigned Integers: Decimal and Binary

- Unsigned integers are always positive:

  ```
  unsigned int i = 12345;
  ```

- To understand binary, recall how decimal numbers "work"

Decimal: Base 10 Example

Each digit adds on a power 10

$$
\begin{aligned}
80,345 =& 5 \times 10^0 + && \text{5 ones} \\
& 4 \times 10^1 + && \text{40 tens} \\
& 3 \times 10^2 + && \text{300 hundreds} \\
& 0 \times 10^3 + && \text{0 thousands} \\
& 8 \times 10^4 && 80,000 \ldots \\
5 + 40 &+ 300 + 80,000
\end{aligned}
$$

Binary: Base 2 Example

Each digit adds on a power 2

$$
\begin{aligned}
11001_2 =& 1 \times 2^0 + && \text{1 ones} \\
& 0 \times 2^1 + && \text{0 twos} \\
& 0 \times 2^2 + && \text{0 fours} \\
& 1 \times 2^3 + && \text{8 eights} \\
& 1 \times 2^4 + && \text{16 sixteens} \\
=& 1 + 8 + 16 && = 25
\end{aligned}
$$

So, $11001_2 = 25_{10}$

Exercise: Convert Binary to Decimal

Base 2 Example:

$$
\begin{aligned}
11001 = & 1 \times 2^0 + & 1 \\
& 0 \times 2^1 + & 0 \\
& 0 \times 2^2 + & 0 \\
& 1 \times 2^3 + & 8 \\
& 1 \times 2^4 + & 16 \\
= & 1 + 8 + 16 & = 25
\end{aligned}
$$

So, $11001_2 = 25_{10}$

Try With a Neighbor

Convert the following two numbers from base 2 (binary) to base 10 (decimal)

- 111
- 11010
- 01100001

Answers: Convert Binary to Decimal

$$
\begin{aligned}
111_2 =& 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 \\
=& 1 \times 4 + 1 \times 2 + 1 \times 1 \\
=& 7_{10} \\
11010_2 =& 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 \\
=& 1 \times 16 + 1 \times 8 + 0 \times 4 + 1 \times 2 + 0 \times 1 \\
=& 26_{10} \\
01100001_2 =& 0 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 \\
& + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\
=& 0 \times 128 + \times 64 + 1 \times 32 + 0 \times 16 \\
& + 0 \times 8 + 0 \times 4 + 0 \times 2 + 1 \times 1 \\
=& 97_{10}
\end{aligned}
$$

Note: last example ignores leading 0's

The Other Direction: Base 10 to Base 2

Converting a number from base 10 to base 2 is easily done using repeated division by 2; keep track of **remainders**

Convert 124 to base 2:

$$
\begin{aligned}
124 \div 2 &= 62 && \text{rem } 0 \\
62 \div 2 &= 31 && \text{rem } 0 \\
31 \div 2 &= 15 && \text{rem } 1 \\
15 \div 2 &= 7 && \text{rem } 1 \\
7 \div 2 &= 3 && \text{rem } 1 \\
3 \div 2 &= 1 && \text{rem } 1 \\
1 \div 2 &= 0 && \text{rem } 1
\end{aligned}
$$

- Last step got 0 quotient so we're done.
- Binary digits are in **remainders in reverse**
- Answer: `1111100`
- Check:
  $0 + 0 + 2^2 + 2^3 + 2^4 + 2^5 + 2^6 = 4 + 8 + 16 + 32 + 64 = 124$

Decimal, Hexadecimal, Octal, Binary

- Numbers exist independent of any writing system
- Can write the same number in a variety of bases
- C provides syntax for most common bases used in computing

	Decimal	Binary	Hexadecimal	Octal
Base	10	2	16	8
Mathematical	125	1111101_2	$7D_{16}$	175_8
C Prefix	None	`0b...`	`0x..`	`0...`
C Example	125	`0b1111101`	`0x7D`	`0175`

- **Hexadecimal** often used to express long-ish byte sequences
  Larger than base 10 so for 10-15 uses letters A-F
- **Examine** `number_writing.c` and `table.c` for patterns
- **Expectation**: Gain familiarity with doing conversions between bases as it will be useful in practice

Hexadecimal: Base 16

- Hex: compact way to write bit sequences
- One byte is 8 bits
- Each Hex character represents 4 bits
- **Each Byte is 2 Hex Digits**

```
|-----------+----------------+-----|
| Byte      | Hex            | Dec |
|-----------+----------------+-----|
| 0101 0111 | 57 = 5*16 + 7  |  87 |
| 5    7    |                |     |
| 0011 1100 | 3C = 3*16 + 12 |  60 |
| 3    C=12 |                |     |
| 1110 0010 | E2 = 14*16 + 2 | 226 |
| E=14 2    |                |     |
|-----------+----------------+-----|
```

Hex to 4 bit equivalence

Dec	Bits	Hex
0	0000	0
1	0001	1
2	0010	2
3	0011	3
4	0100	4
5	0101	5
6	0110	6
7	0111	7
8	1000	8
9	1001	9
10	1010	A
11	1011	B
12	1100	C
13	1101	D
14	1110	E
15	1111	F

Exercise: Conversion Tricks for Hex and Octal

Examples shown in this week's HW, What tricks are illustrated?

```
|---------+--------------+------------------+-----------------------|
| Decimal | Byte = 8bits | Byte by 4        | Hexadecimal           |
|---------+--------------+------------------+-----------------------|
|      87 |     01010111 | bin: 0101 0111   | 57 = 5*16 + 7         |
|         |              | hex: 5    7      | hex  dec              |
|         |              |                  |                       |
|      60 |     00111100 | bin: 0011 1100   | 3C = 3*16 + 12        |
|         |              | hex: 3    C=12   | hex  dec              |
|         |              |                  |                       |
|     226 |     11100010 | bin: 1110 0010   | E2 = 14*16 + 2        |
|         |              | hex: E=14 2      | hex  dec              |
|---------+--------------+------------------+-----------------------|
| Decimal | Byte = 8bits | Byte by 3        | Octal                 |
|---------+--------------+------------------+-----------------------|
|      87 |     01010111 | bin: 01 010 111  | 127 = 1*8^2 + 2*8 + 7 |
|         |              | oct: 1  2   7    | oct   dec             |
|         |              |                  |                       |
|      60 |     00111100 | bin: 00 111 100  | 074 = 0*8^2 + 7*8 + 4 |
|         |              | oct: 0  7   4    | oct   dec             |
|         |              |                  |                       |
|     226 |     11100010 | bin: 11 100 010  | 342 = 3*8^2 + 4*8 + 2 |
|         |              | oct: 3  4   2    | oct   dec             |
|---------+--------------+------------------+-----------------------|
```

Answers: Conversion Tricks for Hex and Octal

- Converting between Binary and Hexadecimal is easiest when grouping bits by 4: each 4 bits corresponds to one hexadecimal digit

```
bin: 0101 0111    bin: 1110 0010
hex: 5    7       hex: E=14 2
```

- Converting between Binary and Octal is easiest when grouping bits by 3: each 3 bits corresponds to one octal digit

```
bin: 01 010 111  bin: 11 100 010
oct: 1  2   7    oct: 3  4   2
```

Character Coding Conventions

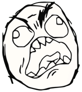
- Would be hard for people to share words if they interpretted bits as letters differently
- **ASCII**: American Standard Code for Information Interchange
  An old standard for bit/character correspondence
- 7 bits per character, includes upper, lower case, punctuation

Dec	Hex	Binary	Char	Dec	Hex	Binary	Char
65	41	01000001	A	78	4E	01001110	N
66	42	01000010	B	79	4F	01001111	O
67	43	01000011	C	80	50	01010000	P
68	44	01000100	D	81	51	01010001	Q
69	45	01000101	E	82	52	01010010	R
70	46	01000110	F	83	53	01010011	S
71	47	01000111	G	84	54	01010100	T
72	48	01001000	H	85	55	01010101	U
73	49	01001001	I	86	56	01010110	V
74	4A	01001010	J	87	57	01010111	W
75	4B	01001011	K	88	58	01011000	X
76	4C	01001100	L	89	59	01011001	Y
77	4D	01001101	M	90	5A	01011010	Z
91	5B	01011101	[	97	61	01100001	a
92	5C	01011110	\	98	62	01100010	b

Unicode

- World: Why can't I write

  컴퓨터

  in my code/web address/email?
- America: ASCII has 128 chars. Deal with it.
- World: Seriously?
- America: We invented computers. 'Merica!
- World: 
- America: ··· Unicode?
- World: But my language takes more bytes than American.
- America: Deal with it. 'Merica!

- ASCII Uses 7 bits per char, limited to 128 characters
- UTF-8 uses **1-4 bytes per character** to represent **many** more characters (1,112,064 *codepoints*)
- Uses 8th bit in a byte to indicate extension to more than a single byte
- Requires software to understand coding convention allowing broader language support
- ASCII is a proper subset of UTF-8 making UTF-8 backwards compatible and increasingly popular

Binary Integer Addition/Subtraction

Adding/subtracting in binary works the same as with decimal EXCEPT that carries occur on values of 2 rather than 10

```
ADDITION #1
   1 11     <-carries
   0100 1010 = 74
 + 0101 1001 = 89
------------
   1010 0011 = 163

ADDITION #2
   1111   1 <-carries
   0110 1101 = 109
 + 0111 1001 = 121
------------
   1110 0110 = 230
```

```
SUBTRACTION #1
          ? <-carries
   0111 1001 = 121
 - 0001 0011 =  19
------------
   VVVVVVVVVVVVV
   VVVVVVVVVVVVV
   VVVVVVVVVVVVV
        x12 <-carries
   0111 0001 = 119
 - 0001 0011 =  19
------------
   0110 0110 = 102
```

Two's Complement Integers: Representing Negative Values

- To represent negative integers, must choose a coding system
- **Two's complement** is the most common for this
- Alternatives exist
  - Signed magnitude: leading bit indicates pos (0) or neg (1)
  - One's complement: invert bits to go between positive negative
- Great advantage of two's complement: **signed and unsigned arithmetic are identical**
- Hardware folks only need to make one set of units for both unsigned and signed arithmetic

Summary of Two's Complement

Short explanation: most significant bit is associated with a negative power of two.

```
UNSIGNED BINARY         TWO's COMPLEMENT (signed)
---------------         -------------------------
7654 3210 : position    7654 3210 : position
ABCD EFGH : 8 bits      ABCD EFGH : 8-bits
A: 0/1 * +(2^7) *POS*   A: 0/1 * -(2^7) *NEG*
B: 0/1 * +(2^6)         B: 0/1 * +(2^6)
C: 0/1 * +(2^5)         C: 0/1 * +(2^5)
...                     ...
H: 0/1 * +(2^0)         H: 0/1 * +(2^0)

UNSIGNED BINARY         TWO's COMPLEMENT (signed)
---------               ---------
7654 3210 : position    7654 3210 : position
1000 0000 = +128        1000 0000 = -128
1000 0001 = +129        1000 0001 = -127 = -128+1
1000 0011 = +131        1000 0011 = -125 = -128+1+2
1111 1111 = +255        1111 1111 = -1   = -128+1+2+4+..+64
0000 0000 = 0           0000 0000 =  0         [   +127    ]
0000 0001 = +1          0000 0001 = +1
0000 0101 = +5          0000 0101 = +5
0111 1111 = +127        0111 1111 = +127
```

Two's Complement Notes

- Leading 1 indicates negative, 0 indicates positive
- All 0's = Zero
- Positive numbers are identical to unsigned

Conversion Trick

Positive → Negative

- **Invert bits, Add 1**

Negative → Positive

- **Invert bits, Add 1**

Same trick works both ways, implemented in hardware for the **unary minus** operator as in

```
int y = -x;
```

```
~ 0110 1000   +104 : negate
-----------
  1001 0111   inverted
+         1
-----------
  1001 1000 = -104

~ 1001 1000 = -104 : negate
-----------
  0110 0111 = +103 inverted
+         1
-----------
  0110 1000 = +104
```

Add Pos/Neg should give 0

```
1 1111     <-carries
  0110 1000 = +104
+ 1001 1000 = -104
-----------
x 0000 0000 = zero
```

Overflow

- Sums that exceed the representation of the bits associated with the integral type **overflow**
- Excess significant bits are **dropped**
- Addition can result in a sum smaller than the summands, even for two positive numbers (!?)
- Integer arithmetic in fixed bits is a mathematical **ring**

Examples of Overflow in 8 bits

```
ADDITION #3 OVERFLOW    ADDITION #4 OVERFLOW
 1 1111 111 <-carries    1         1 <-carries
   1111 1111 = 255         1010 1001 = 169
 + 0000 0001 =   1       + 1100 0001 = 193
------------            ------------
 1 0000 0000 = 256       1 0110 1010 = 362
 x drop 9th bit          x drop 9th bit
------------            ------------
   0000 0000 = 0           0110 1010 = 106
```

Underflow

- **Underflow** occurs in unsigned arithmetic when values go below 0 (no longer positive)
- Pretend that there is an extra significant bit to carry out subtraction
- Subtracting a positive integer from a positive integer may result in a **larger** positive integer (?!?)
- Integer arithmetic in fixed bits is a mathematical **ring**

Examples of 8-bit Underflow

```
SUBTRACTIION #2 UNDERFLOW
          ?<-carries
   0000 0000 =   0
 - 0000 0001 =   1
------------
VVVVVVVVVVVVV
          ?<-carries
 1 0000 0000 = 256 (pretend)
 - 0000 0001 =   1
------------
VVVVVVVVVVVVV
 x        2<-carries
 0 1111 1110 = 256
 - 0000 0001 =   1
------------
   1111 1111 = 255
```

Overflow and Underflow In C Programs

- See `over_under_flow.c` for demonstrations in a C program.
- **No runtime errors** for under/overflow
- Good for hashing and cryptography
- Bad for most other applications: system critical operations should use checks for over-/under-flow
- See textbook Ariane Rocket Crash which was due to overflow of an integer converted from a floating point value
- At the assembly level, there are condition codes indicating that overflow has occurred

Endinaness: Byte ordering in Memory

- Single bytes like ASCII characters lay out sequentially in memory in increasing address
- Multi-byte entities like 4-byte `ints` require decisions on byte ordering
- We think of a 32-bit int like this

```
Binary:   0000 0000  0000 0000  0001 1000  1110 1001
          0    0     0    0     1    8     E    9
Hex    :  000018E9
Decimal:  6377
```

- But need to assign memory addresses to each byte
  - Little Endian: least significant byte early
  - Big Endian: most significant byte early
- Example: Integer starts at address #1024

```
           Address
LittleEnd: #1027      #1026      #1025      #1024
Binary:    0000 0000  0000 0000  0001 1000  1110 1001
           0    0     0    0     1    8     E    9
BigEnd:    #1024      #1025      #1026      #1027
           Address
```

Little Endian vs. Big Endian

- Most modern machines use **little endian** by default
- Processor may actually support big endian
- Both Big and Little Endian have engineering trade-offs
- At one time debated hotly among hardware folks: a la Gulliver's Travels conflicts
- **Intel Chips** were little endian and have dominated computing for several decades, set the precedent for modern platforms
- Big endian byte order shows up in **network programming**: sending bytes over the network is done in big endian ordering
- **Examine** `show_endianness.c` : uses C code to print bytes in order, reveals whether a machine is Little or Big Endian

Output of `show_endianness.c`

```
> cat show_endianness.c
// Show endiannes layout of a binary number in memory. Most machines
// are little endian so bytes will print leas signficant earlier.
#include <stdio.h>

int main(){
  int bin = 0b00000000000000000001100011101001;    // 6377
  //          |   |   |   |   |   |   |   |
  //          0   0   0   0   1   8   e   9
  printf("%d\n%x\n",bin,bin);                      // show decimal/hex of binary
  unsigned char *ptr = (unsigned char *) &bin;    // pointer to beginning of bin
  for(int i=0; i<4; i++){                          // print bytes of bin from low to high
    printf("%hhx ", ptr[i]);                       // memory address
  }                                                // '%hhx' : 1-byte char in hex
  printf("\n");                                    // '%hx'  : 2-byte short in hex
  return 0;                                        // '%x'   : 4-byte int in hex
}

> gcc show_endianness.c
> ./a.out
6377
18e9
e9 18 0 0
```

Notice: `num` prints with value `18e9` but bytes appear in reverse order `e9 18` when looking at memory

Integer Ops and Speed

- Along with Addition and Subtraction, **Multiplication and Division** can also be done in binary
- Algorithms are the same as base 10 but more painful to do by hand
- This pain is reflected in hardware speed of these operations
- The **Arithmetic and Logic Unit (ALU)** does integer ops in the machine
- A **clock** ticks in the machine at some rate like 3Ghz (3 billion times per second)
- Under ideal circumstances, typical ALU Op speeds are

Operation	Cycles
Addition	1
Logical	1
Shifts	1
Subtraction	1
Multiplication	3
Division	>30

- Due to disparity, it is worth knowing about **relation** between multiply/divide and **bitwise** operations
- Compiler often uses such tricks: shift rather than multiply/divide

Mangling bits puts hair on your chest

Below illustrates difference between logical and bitwise operations.

```
int xl = 12 || 10; // truthy (Logical OR)
int xb = 12 |  10; // 14     (Bitwise OR)
int yl = 12 && 10; // truthy (Logical AND)
int yb = 12 &  10; // 8      (Bitwise AND)
int zb = 12 ^  10; // 6      (Bitwise XOR)
int wl = !12;      // falsey (Logical NOT)
int wb = ~12;      // 3      (Bitwise NOT/INVERT)
```

- Bitwise ops evaluate on a per-bit level
- 32 bits for `int`, 4 bits shown

```
Bitwise OR    Bitwise AND    Bitwise XOR   Bitwise NOT
  1100 = 12     1100 = 12      1100 = 12
| 1010 = 10   & 1010 = 10    ^ 1010 = 10   ~ 1100 = 12
-----------   -----------    -----------   -----------
  1110 = 14     1000 =  8      0110 =  6     0011 =  3
```

Bitwise Shifts

- **Shift** operations move bits within a field of bits
- Shift operations are

```
x = y << k; // left  shift y by k bits, store in x
x = y >> k; // right shift y by k bits, store in x
```

- All integral types can use shifts: `long`, `int`, `short`, `char`
- **Not applicable** to pointers or floating point
- Examples in 8 bits

```
//         76543210
char x = 0b00010111; // 23
char y = x << 2;     // left shift by 2
//   y = 0b01011100; // 92
//   x = 0b00010111; // not changed
char z = x >> 3;     // right shift by 3
//   z = 0b00000010; // 2
//   x = 0b00010111; // not changed
char n = 0b10000000; // -128, signed
char s = n >> 4;     // right shift by 4
//   s = 0b11111000; // -8, sign extension
//   right shift >> is "arithmetic"
```

Shifty Arithmetic Tricks

- Shifts with add/subtract can be used instead of multiplication and division
- Turn on optimization: `gcc -O3 code.c`
- Compiler automatically does this if it thinks it will save cycles
- *Sometimes* programmers should do this but better to convince compiler to do it for you, **comment** if doing manually

Multiplication

```
//           76543210
char  x = 0b00001010;   // 10
char x2 = x << 1;       // 10*2
//   x2 = 0b00010100;   // 20
char x4 = x << 2;       // 10*4
//   x4 = 0b00101000;   // 40
char x7 = (x << 3)-x;   // 10*7
//   x7 = (x * 8)-x;    // 10*7
//   x7 = 0b01000110;   // 70
//           76543210
```

Division

```
//           76543210
char y  = 0b01101110;  // 110
char y2 = y >> 1;      // 110/2
//   y2 = 0b00110111;  // 55
char y4 = y >> 2;      // 110/4
//   y4 = 0b00011011;  // 27
char z  = 0b10101100;  // -84
char z2 = z >> 2;      // -84/4
//   z2 = 0b11101011;  // -21
//   right shift sign extension
```

Exercise: Checking / Setting Bits

Use a combination of bit shift / bitwise logic operations to⋯

1. Check if bit `i` of `int x` is set (has value 1)
2. Clear bit `i` (set bit at index `i` to value 0)

Show C code for this

```
{
  int x = ...;
  int i = ...;
  if( ??? ) { // ith bit of x is set
    printf("set!\n");
  }

  i = ...;
  ???;
  printf("ith bith of x now cleared to 0\n");
}
```

Answers: Checking / Setting Bits

1. Check if bit `i` of `int x` is set (has value 1)

```
int x = ...;
int mask = 1; // or 0b0001 or 0x01 ...
int shifted  = mask << i; // shifted   0b00...010..00
if(x & shifted){          //       x & 0b10...010..01
  ...                     //       ------------------
}                         //           0b00...010..00
```

2. Clear bit `i` (set bit at index `i` to value 0)

```
int x = ...;
int mask = 1; // or 0b0001 or 0x01 ...
int shifted    = mask << i;  // shifted   0b00...010..00
int inverted   = ~shifted;   // inverted  0b11...101..11
x = x & inverted;            //       x & 0b10...010..01
  ...                        //       ------------------
                             //           0b10...000..01
```

Showing Bits

- `printf()` capabilities:

%d	as Decimal
%x	as Hexadecimal
%o	as Octal
%c	as Character

- No specifier for binary
- Can construct such with bitwise operations
- Code pack contains two codes to do this
  - `printbits.c`: single args printed as 32 bits
  - `showbits.c`: multiple args printed in binary, hex, decimal

- Showing bits usually involves shifting and bitwise AND &
- Example from `showbits.c`

```
#define INT_BITS 32

// print bits for x to screen
void showbits(int x){
  for(int i=INT_BITS-1; i>=0; i--){
    int mask = 1 << i;
    if(mask & x){
      printf("1");
    } else {
      printf("0");
    }
  }
}
```

Bit Masking

- Semi-common for functions to accept bit patterns which indicate true/false options
- Frequently makes use of bit **masks** which are constants associated with specific bits
- Example from earlier: Unix permissions might be···

```
#define S_IRUSR 0b100000000 // User   Read
#define S_IWUSR 0b010000000 // User   Write
#define S_IXUSR 0b001000000 // User   Execute
#define S_IRGRP 0b000100000 // Group  Read
...
#define S_IWOTH 0b000000010 // Others Write
#define S_IXOTH 0b000000001 // Others Execute
```

- Use them to create options to C functions like

```
int permissions = S_IRUSR|S_IWUSR|S_RGRP;
chmod("/home/kauffman/solution.zip",permissions);
```

Unix Permissions with Octal

- Octal arises associated with **Unix file permissions**
- Every file has 3 permissions for 3 entities
- Permissions are true/false so a single bit will suffice

- `ls -l`: long list files, shows permissions
- `chmod 665 somefile.txt`: change permissions of `somefile.txt` to those shown to the right
- `chmod 777 x.txt`: read / write / exec for everyone
- chmod also honors letter versions like `r` and `w`
- `chmod u+x script.sh # make file executable`

```
binary      octal
110110101 = 665
rw-rw-r-x somefile.txt
U  G  O
S  R  T
E  O  H
R  U  E
   P  R

Readable chmod version:
chmod u=rw,g=rw,o=rx somefile.txt
```